

CMSC 3540I: The Interplay of Economics and ML  
(Winter 2024)

# The Value and Pricing of Information

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Instructor: Haifeng Xu



# Outline

- Bayesian Persuasion and Information Selling
- Sell to a Single Decision Maker
- Sell to Multiple Decision Makers

# Motivation: Selling Information

- Car/house inspections



- Financial advices



- Credit report



- Consumer data



# Motivation: Selling Information

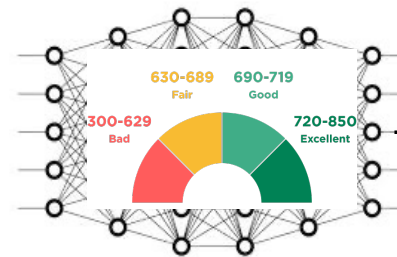
- Car/house inspections



- Financial advices



- Credit report



→ Prob. of default

- Consumer data



→ Prob. of purchase/conversion

# Persuasion vs Information Selling

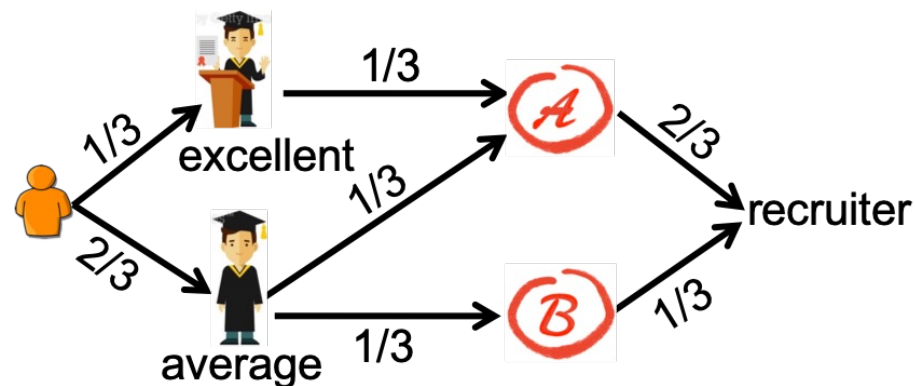
- In persuasion, we selectively reveal information to induce actions that we like



When selling information, we reveal information for a profit

# Recap: Model of Bayesian Persuasion

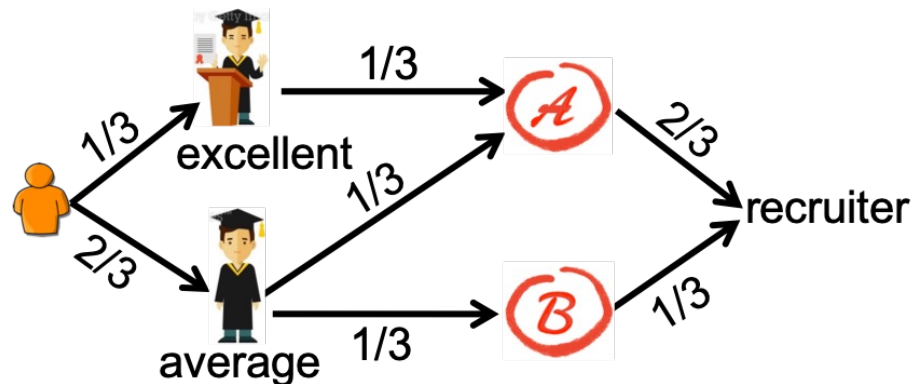
- Two players: persuader (**Sender, she**), decision maker (**Receiver he**)
  - Example: advisor = sender, recruiter = receiver
- Receiver looks to take an action  $i \in [n] = \{1, 2, \dots, n\}$ 
  - Receiver utility  $r(i, \theta)$       $\theta \in \Theta$  is a random **state of nature**
  - Sender utility  $s(i, \theta)$
- Both players know  $\theta \sim$  *prior dist.*  $\mu$ , but Sender has an **informational advantage** – she can observe realization of  $\theta$
- Sender reveal partial information via a signaling scheme



# (Simplified) Model of Selling Information

seller

- Two players: ~~persuader~~ (Sender, she), decision maker (Receiver he)
  - Example: advisor = sender, recruiter = receiver
- Receiver looks to take an action  $i \in [n] = \{1, 2, \dots, n\}$ 
  - Receiver utility  $r(i, \theta)$       $\theta \in \Theta$  is a random state of nature
  - Sender utility  ~~$s(i, \theta)$~~  — payment from the receiver
- Both players know  $\theta \sim \text{prior dist. } \mu$ , but Sender has an informational advantage – she can observe realization of  $\theta$
- Sender reveal partial information via a signaling scheme



# How to Sell Information Optimally?

➤ For any signaling scheme, seller knows how much it improves buyer's expected utility

- The value of any signaling scheme is known

1. Receiver utility under **no information**:  $\max_i \sum_{\theta \in \Theta} r(i, \theta) \cdot \mu(\theta)$

2. Receiver utility under **any  $\pi$** :  $\sum_{\sigma} \Pr(\sigma) \cdot R(\sigma)$

$$\text{where } R(\sigma) = \max_{i \in [n]} \left[ \sum_{\theta \in \Theta} r(i, \theta) \cdot \frac{\pi(\sigma, \theta) \cdot \mu(\theta)}{\Pr(\sigma)} \right]$$

➤ How to maximize revenue?

- Reveal full information helps the buyer the most. Why?
- So OPT is to charge him following amount and **then** reveal  $\theta$  directly

$$\text{Payment} = \sum_{\theta \in \Theta} \mu(\theta) \cdot \left[ \max_i u(i, \theta) \right] - \max_i \sum_{\theta \in \Theta} \mu(\theta) \cdot u(i, \theta)$$



Buyer expected utility if learns  $\theta$  precisely



# How to Sell Information Optimally?

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**Q:** Are we done?

No – in pricing problems, we typically do not know how much buyer values our “product”

# Outline

- Bayesian Persuasion and Information Selling
- Sell to a Single Decision Maker
- Sell to Multiple Decision Makers

# (True) Model of Selling Information

- Sender = seller, Receiver = buyer who is a decision maker
- Buyer takes an action  $i \in [n] = \{1, \dots, n\}$
- Buyer has a utility function  $u(i, \theta; t)$  where
  - $\theta \sim \text{dist. } \mu$  is a random state of nature
  - $t \sim \text{dist. } f$  captures buyer's (private) utility type

## Remarks:

- $u, \mu, f$  are public knowledge
- Assume  $\theta, t$  are independent
- Seller observes  $\theta$  but does not know buyer's type  $t$
- Buyer knows his own type  $t$  but does not know  $\theta$

# Key Challenge

The class of mechanisms is too broad

- The mechanism will: (1) elicit private info from buyer; (2) reveal info based on realized  $\theta$ ; (3) charge buyer
- May interact with buyer for many rounds
- Buyer may misreport his private type  $t$

# Key Challenge

The class of mechanisms is too broad

. . . but, at the end of the day, the buyer of type  $t$  is charged some amount  $x_t$  in expectation and learns a posterior belief about  $\theta$

**Theorem (Revelation Principle).** Any information selling mechanism is “equivalent” to a **direct and truthful revelation mechanism**:

1. Ask buyer to report type  $t$
2. Charge buyer  $x_t$  and reveal info to buyer via signaling scheme  $\pi_t$  that use  $n$  signals (as action recommendations)

Moreover, the mechanism is incentive compatible (IC) – it is the buyer’s best interest to truthfully report  $t$

- Optimal mechanism reduces to computing an IC menu  $\{x_t, \pi_t\}_t$
- Proof omitted here

# The Optimal Mechanism

## The Consulting Mechanism [CXZ, SODA'20]

1. Elicit buyer type  $t$
2. Charge buyer  $x_t$
3. Observe realized state  $\theta$  and recommend action  $i$  to the buyer with probability  $\pi_t(\sigma_i, \theta)$

- Will be incentive compatible – reporting true  $t$  is optimal
- The recommended action is guaranteed to be the optimal action for buyer  $t$  given his information
- $\{x_t, \pi_t\}_t$  is public knowledge, and computed by LP

**Theorem.** Consulting mechanism with  $\{x_t, \pi_t\}_t$  computed by the following program is the optimal mechanism.

# Computing the Optimal Mechanism

Optimal  $\{x_t, \pi_t\}_t$  can be computed by a convex program

- Variables:  $\pi_t(\sigma_i, \theta)$  = prob of sending  $\sigma_i$  conditioned on  $\theta$  for each  $t$
- Variable  $x_t$  is the payment from buyer type  $t$

Expected revenue

$$\max \quad \sum_t f(t) \cdot x_t$$

$$\begin{aligned} \text{s.t.} \quad & \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \\ & \geq \sum_i \max_j \left[ \sum_{\theta} \mu(\theta) \pi_{t'}(\sigma_i, \theta) u(j, \theta; t) \right] - x_{t'}, && \text{for } t' \neq t \\ & \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \max_i \sum_{\theta} \mu(\theta) u(i, \theta; t), && \text{for } t \\ & \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \geq \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(j, \theta; t), && \text{for } i \neq j, t \\ & \sum_i \pi_t(\sigma_i, \theta) = 1, && \text{for } \theta, t \\ & \pi_t(\sigma_i, \theta) \geq 0, && \text{for } t, \sigma_i, \theta \end{aligned}$$

# Computing the Optimal Mechanism

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- Variables:  $\pi_t(\sigma_i, \theta)$  = prob of sending  $\sigma_i$  conditioned on  $\theta$  for each  $t$
- Variable  $x_t$  is the payment from buyer type  $t$

Truthfully reporting true  $t$  is optimal

$$\max \sum_t f(t) \cdot x_t$$

$$\text{s.t. } \begin{cases} \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \\ \geq \sum_i \max_j \left[ \sum_{\theta} \mu(\theta) \pi_{t'}(\sigma_i, \theta) u(j, \theta; t) \right] - x_{t'}, & \text{for } t' \neq t \\ \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \max_i \sum_{\theta} \mu(\theta) u(i, \theta; t), & \text{for } t \\ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \geq \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(j, \theta; t), & \text{for } i \neq j, t \\ \sum_i \pi_t(\sigma_i, \theta) = 1, & \text{for } \theta, t \\ \pi_t(\sigma_i, \theta) \geq 0, & \text{for } t, \sigma_i, \theta \end{cases}$$



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- Variable  $x_t$  is the payment from buyer type  $t$

Participation is no worse than not

$$\max \quad \sum_t f(t) \cdot x_t$$

$$\text{s.t.} \quad \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \sum_i \max_j \left[ \sum_{\theta} \mu(\theta) \pi_{t'}(\sigma_i, \theta) u(j, \theta; t) \right] - x_{t'}, \quad \text{for } t' \neq t$$

$$\sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \max_i \sum_{\theta} \mu(\theta) u(i, \theta; t), \quad \text{for } t$$

$$\sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \geq \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(j, \theta; t), \quad \text{for } i \neq j, t$$

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# Computing the Optimal Mechanism

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$$\begin{aligned}
 \max \quad & \sum_t f(t) \cdot x_t \\
 \text{s.t.} \quad & \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \\
 & \geq \sum_i \max_j \left[ \sum_{\theta} \mu(\theta) \pi_{t'}(\sigma_i, \theta) u(j, \theta; t) \right] - x_{t'}, \quad \text{for } t' \neq t \\
 & \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \max_i \sum_{\theta} \mu(\theta) u(i, \theta; t), \quad \text{for } t \\
 & \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \geq \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(j, \theta; t), \quad \text{for } i \neq j, t \\
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 \end{aligned}$$

Similar to constraints in persuasion

# Computing the Optimal Mechanism

Optimal  $\{x_t, \pi_t\}_t$  can be computed by a convex program

- Variables:  $\pi_t(\sigma_i, \theta)$  = prob of sending  $\sigma_i$  conditioned on  $\theta$  for each  $t$
- Variable  $x_t$  is the payment from buyer type  $t$

- A convex function of variables
- Can be converted to an LP

$$\max \sum_t f(t) \cdot x_t$$

$$\text{s.t.} \quad \sum_i \left[ \sum_{\theta} \mu(\theta) \pi_t(\sigma_i, \theta) u(i, \theta; t) \right] - x_t \geq \sum_i \max_j \left[ \sum_{\theta} \mu(\theta) \pi_{t'}(\sigma_i, \theta) u(j, \theta; t) \right] - x_{t'}, \quad \text{for } t' \neq t$$

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$$\pi_t(\sigma_i, \theta) \geq 0, \quad \text{for } t, \sigma_i, \theta$$

# Practical Mechanisms?

What the mechanism is like?

- Generally, the optimal solution to the previous LP has no structure neither any interpretation
- Nevertheless, closed-form optimal solution is possible for more structured problems

# Recall Model of Selling Information

- Sender = seller, Receiver = buyer who is a decision maker (DM)
- Buyer takes an action  $i \in [n] = \{1, \dots, n\}$
- Buyer has a utility function  $u(i, q; t)$  where
  - $q \sim \text{dist. } \mu$  is a random state of nature
  - $t \sim \text{dist. } f$  captures buyer's (private) utility type

Remarks:

- $u, \mu, f$  are public knowledge
- Assume  $q, t$  are independent

# Selling Information to a Binary DM

- Sender = seller, Receiver = buyer who is a decision maker (DM)
- Buyer takes an action  $i \in \{0,1\}$ : an **active action 1** and a **passive action 0**
  - Active action: approve loan, buy a car, invest stock X, etc.
- Buyer has a utility function  $u(i, q; t)$  where  $\begin{cases} u(0, q; t) \equiv 0 \\ u(1, q; t) = v(q, t) \end{cases}$ 
  - $q \sim \text{dist. } \mu$  is a random state of nature
  - $t \sim \text{dist. } f$  captures buyer's (private) utility type
- Further assume  $v(q, t)$  is linear and non-decreasing in  $t$

Remarks:

- $u, \mu, f$  are public knowledge
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- Further assume  $v(q, t)$  is linear and non-decreasing in  $t$ 

That is:  $v(q, t) = v_1(q)[t + \rho(q)]$  for some  $v_1(q) \geq 0$

What is the optimal mechanism for this more structured problem?

# An Example



- Buyer is a loan company; action is to approve a loan or not
  - If not approving (action 0), payoff is 0
  - If approving (action 1), payoff is

$$v(q, t) = (1 - q) \times t - 2 \longrightarrow \text{operation cost}$$

$q \in [0,1]$       Revenue

default probability



# Threshold experiments turn out to suffice

Recall  $v(q, t) = v_1(q)[t + \rho(q)]$   
( $q$  is the state unknown to buyer)

**Def.**  $\pi_t$  is a threshold experiment if  $\pi_t$  simply **reveals**  $\rho(q) \geq \theta(t)$  **or not** for some buyer-type-dependent threshold  $\theta(t)$

➤ Threshold is on  $\rho(q)$

# The Magical “Virtual Value Functions”

- Virtual value function turns out to naturally arise at optimal mechanism [Myerson'81]

**Def.** Lower virtual value function:  $\underline{\phi}(t) = t - \frac{1-F(t)}{f(t)}$

# The Magical “Virtual Value Functions”

- Virtual value function turns out to naturally arise at optimal mechanism [Myerson’81]

**Def.** **Lower** virtual value function:  $\underline{\phi}(t) = t - \frac{1-F(t)}{f(t)}$

**Upper** virtual value function:  $\bar{\phi}(t) = t + \frac{F(t)}{f(t)}$

**Mixed** virtual value function:  $\phi_c(t) = c\underline{\phi}(t) + (1 - c)\bar{\phi}(t)$

Note: “upper” or “lower” is due to

$$\underline{\phi}(t) \leq t \leq \bar{\phi}(t)$$

# The Magical “Virtual Value Functions”

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**Def.** **Lower** virtual value function:  $\underline{\phi}(t) = t - \frac{1-F(t)}{f(t)}$

**Upper** virtual value function:  $\bar{\phi}(t) = t + \frac{F(t)}{f(t)}$

**Mixed** virtual value function:  $\phi_c(t) = c\underline{\phi}(t) + (1-c)\bar{\phi}(t)$

- Will assume the virtual value function  $\phi(t)$  is monotone (weakly) increasing in  $t$  (known as the **regularity** assumption)
  - Not crucial – if not monotone, there is a standard procedure to adjust it to make it monotone

# The Optimal Mechanism

**Theorem (Informal, see rigorous statement in [LSX, EC'21]).**

The mechanism with threshold experiments  $\theta^*(t) = -\phi_c^+(t)$  and following payment function represents an optimal mechanism:

$$p^*(t) = \int_{q \in Q} \pi^*(q, t) \mu(q) v(q, t) dq - \int_{t_1}^t \int_{q \in Q} \pi^*(q, x) \mu(q) v_1(q) dq dx$$

where **constant  $c$**  is chosen such that

$$\int_{t_1}^{t_2} \int_{q: \rho(q) \geq \phi_c^+(x)} \mu(q) v_1(q) dq dx = \bar{v}(t_2)$$

# Remarks

- Threshold mechanisms are common in real life
  - House/car inspections, stock recommendations: information seller only need to reveal it “passed” or “deserves a buy” or not
- Optimal mechanism has **personalized** thresholds and payments, tailored to accommodate different level of risk each buyer type can take
  - Different from optimal pricing of physical goods



# Remarks

What if seller is restricted to sell the same information to every buyer? How will revenue change?

- Revenue can be arbitrarily worse
- $1/e$ -approximation of optimal revenue if the *value of full information* as a function of  $t$  is “heavy tail”

# Outline

- Bayesian Persuasion and Information Selling
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# Challenges

- For single decision maker, more information always helps
  - Recall in persuasion, receiver always benefits from signaling scheme
- A fundamental challenge for selling to multiple buyers is that information does not necessarily help them

# Example: More Information Hurts Buyers

- Insurance industry: *insurance company* and *customer*
  - Both are potential information buyers
- Two types of customers: **Healthy** and **Unhealthy**
  - Publicly know,  $\Pr(\text{Healthy}) = 0.9$
- Seller is an information holder, who knows whether any customer is healthy or not

		Insurance company	
		Sell	Not Sell
customer	Buy	(-10, 10)	(-0, 0)
	Not Buy	(0, 0)	(0, 0)

Healthy customer

		Insurance company	
		Sell	Not Sell
customer	Buy	(-10, -50)	(-110, 0)
	Not Buy	(-111, 0)	(-111, 0)

Unhealthy customer

# Example: More Information Hurts Buyers

		Insurance company	
customer		Sell	Not Sell
	Buy	(-10, 10)	(-0, 0)
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Healthy customer, prob = 0.9

		Insurance company	
customer		Sell	Not Sell
	Buy	(-10, -50)	(-110, 0)
	Not Buy	(-111, 0)	(-111, 0)

Unhealthy customer

**Q:** What happens without seller's information ?

- Customer and insurance company will look at expectation
  - Dominant strategy equilibrium is (Buy, Sell)

	Sell	Not Sell
Buy	(-10, 4)	(-11, 0)
Not Buy	(-11.1, 0)	(-11.1, 0)

## Example: More Information Hurts Buyers

		Insurance company	
customer		Sell	Not Sell
	Buy	(-10, 10)	(-0, 0)
	Not Buy	(0, 0)	(0, 0)

Healthy customer, prob = 0.9

		Insurance company	
customer		Sell	Not Sell
	Buy	(-10, -50)	(-110, 0)
	Not Buy	(-111, 0)	(-111, 0)

Unhealthy customer

**Q:** What if seller tells (even only) customer her health status ?

- If Healthy, customer will not buy → utility (0,0) for both
- If Unhealthy, customer will buy → Will not sell, utility (-110,0)
- Customer's reaction reveals his healthy status
- In expectation (-11, 0), and no insurance was sold ever

Recall previous utilities (-10,4)

# Example: More Information Hurts Buyers

Insurance company

customer		Sell	Not Sell
	Buy	(-10, 10)	(-0, 0)
	Not Buy	(0, 0)	(0, 0)

Healthy customer, prob = 0.9

Insurance company

	Sell	Not Sell
Buy	(-10, -50)	(-110, 0)
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Unhealthy customer

**Q:** What if seller tells (even only) customer her health status ?

## Lessons Learned

- Existence of insurance is due to ignorance to our health condition
- Such ignorance benefits both us and insurance companies

# Thank You

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