

Announcement

- Grades for HW2 and project proposal are released

CS6501: Topics in Learning and Game Theory (Fall 2019)

Learning from Strategically Transformed Samples

Instructor: Haifeng Xu

Outline

➤ Introduction

➤ The Model and Results

Signaling

Q: Why attending good universities?

Q: Why publishing and presenting at top conferences?

Q: Why doing internships?

Signaling

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- All in all, these are just **signals** (directly observable) to indicate “excellence” (not directly observable)

Signaling

Q: Why attending good universities?

Q: Why publishing and presenting at top conferences?

Q: Why doing internships?

- All in all, these are just **signals** (directly observable) to indicate “excellence” (not directly observable)
- Asymmetric information between employees and employers

JOB MARKET SIGNALING *

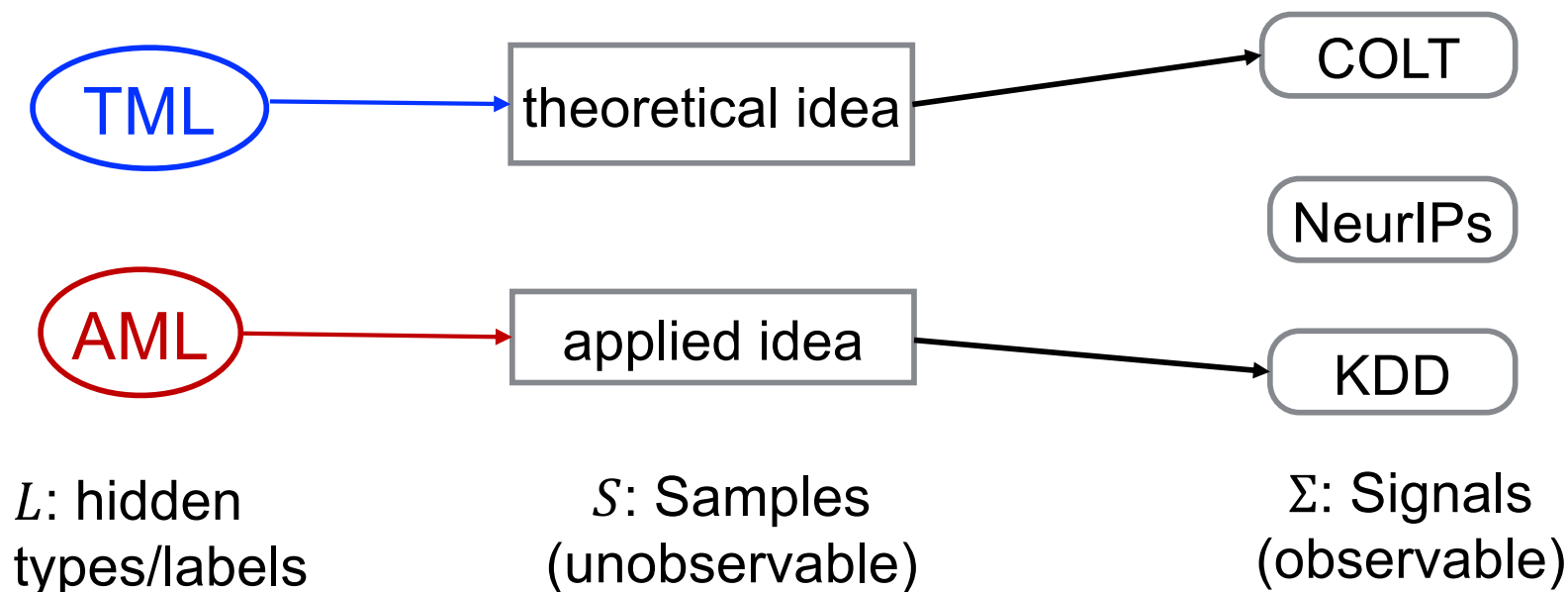
MICHAEL SPENCE

2001 Nobel Econ Price is awarded to research on asymmetric information

Signaling

➤ A simple example

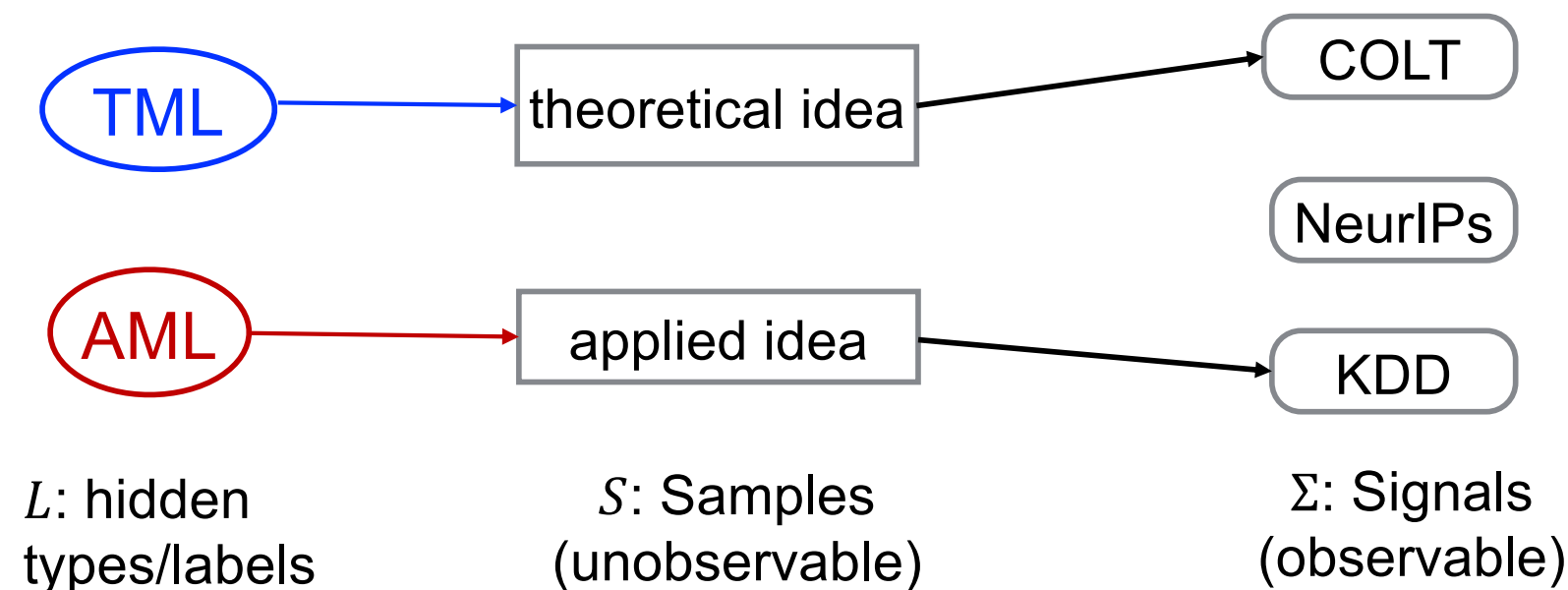
- We want to hire an **Applied ML** researcher
- Only two types of ML researchers in this world
- Easy to tell



Signaling

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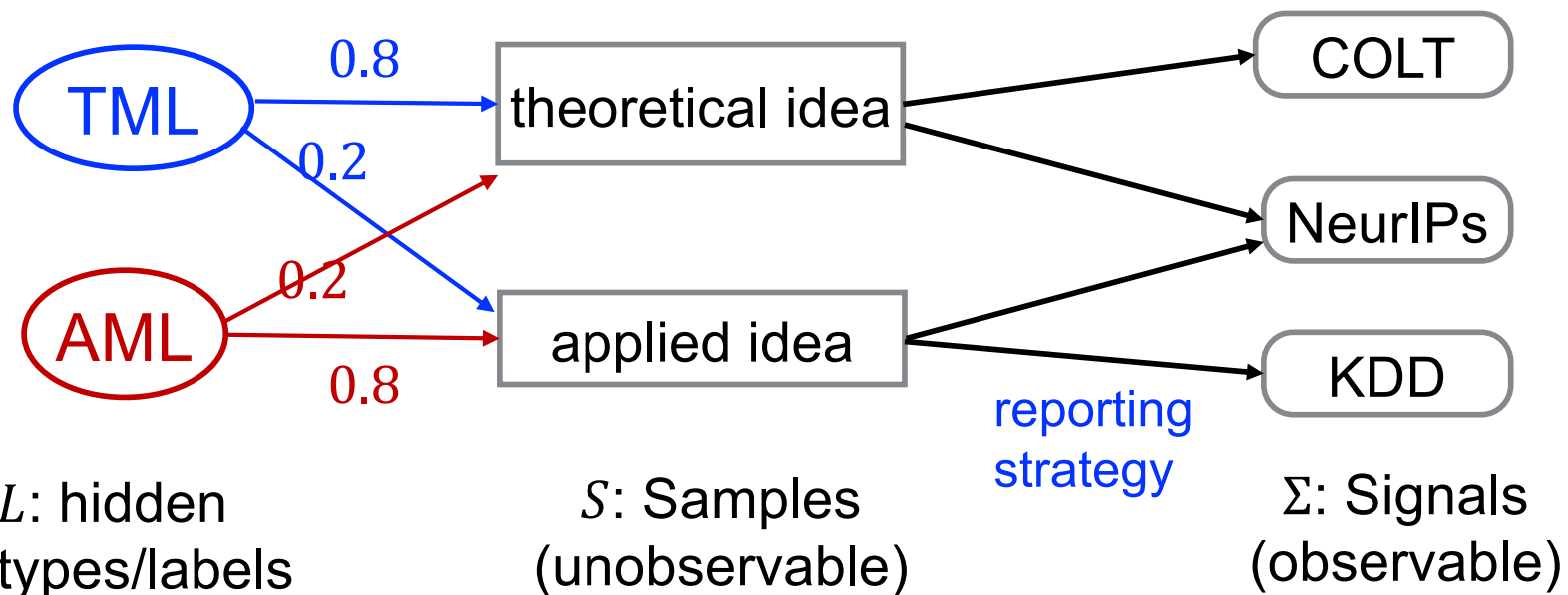


Our world is known to be noisy....

Signaling

➤ A simple example

- We want to hire an **Applied ML** researcher
- Only two types of ML researchers in this world



L : hidden
types/labels

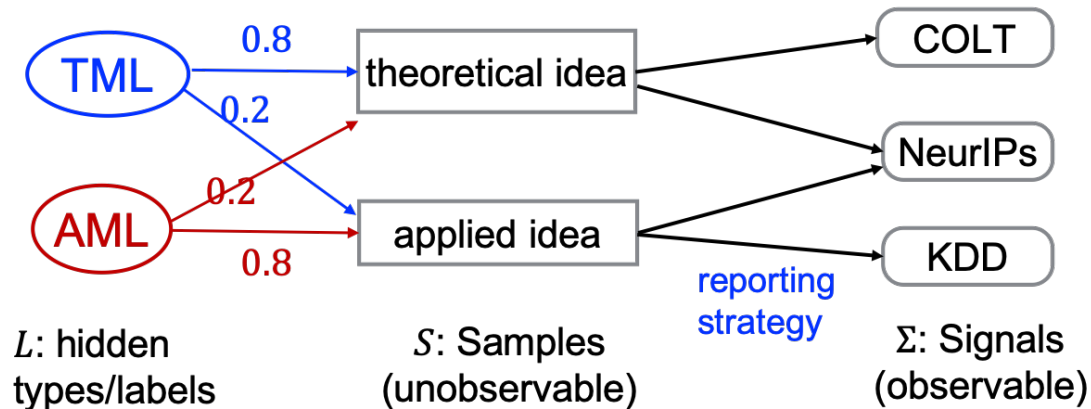
S : Samples
(unobservable)

Σ : Signals
(observable)

$l \in L$ is a distribution
over ideas

generated by l

Signaling



➤ Agent's problem:

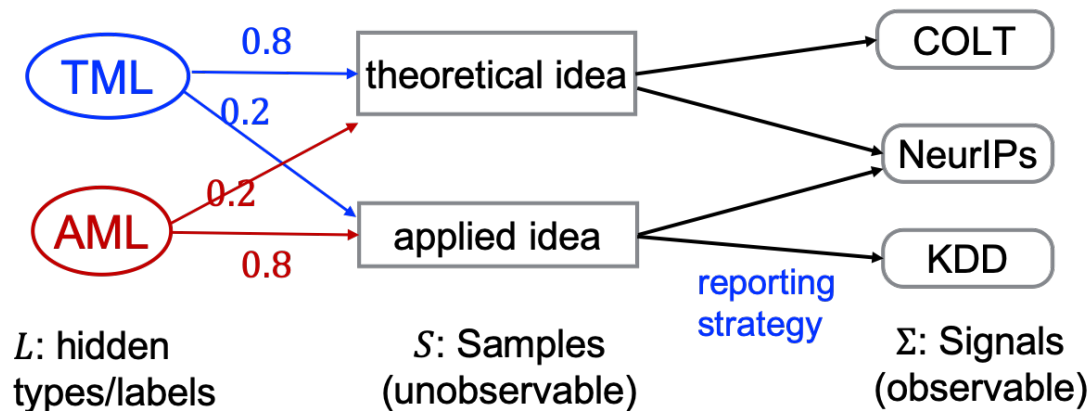
- How do I distinguish myself from other types?
- How many ideas do I need for that?

➤ Principle's problem:

- How do I tell AML agents from others (a classification problem)?
- How many papers should I expect to read?

Answers for this particular instance?

Signaling



➤ Agent's problem:

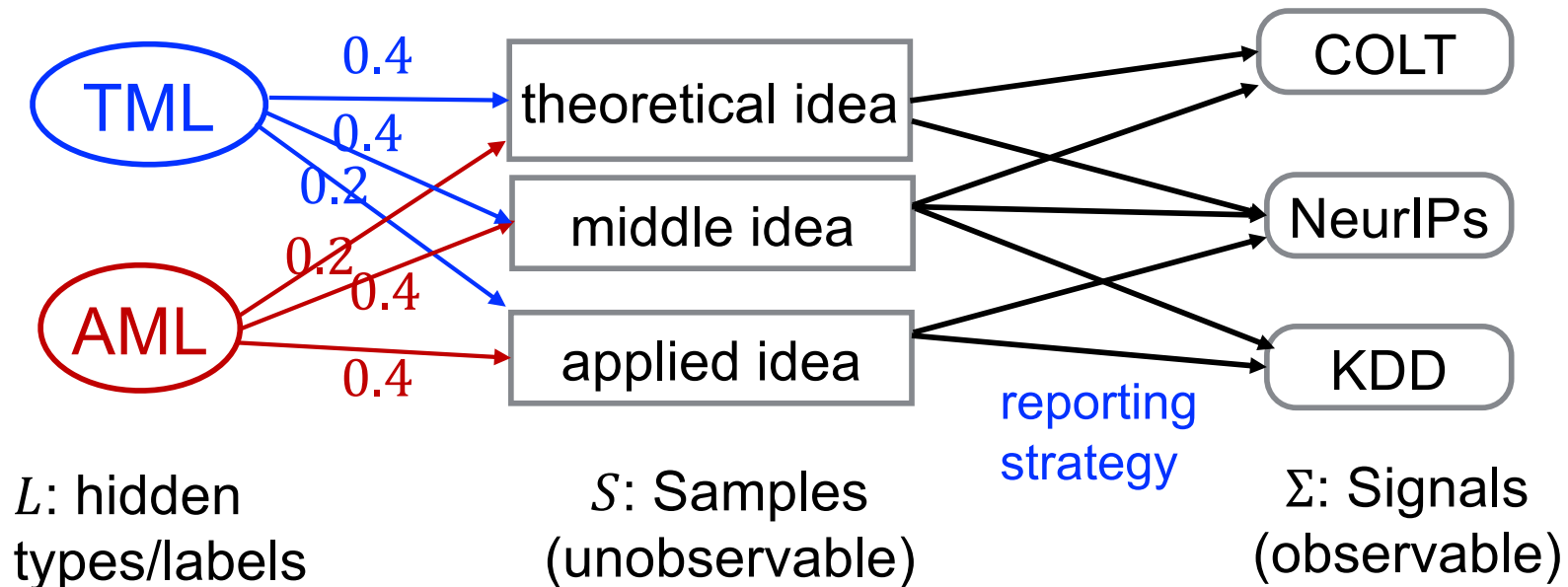
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➤ Principle's problem:

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Generally, classification with strategically transformed samples

What Instances May Be Difficult?



Intuitions

- Agent: try to report as far from others as possible
- Principal: examine a set of signals that maximally separate AML from TML

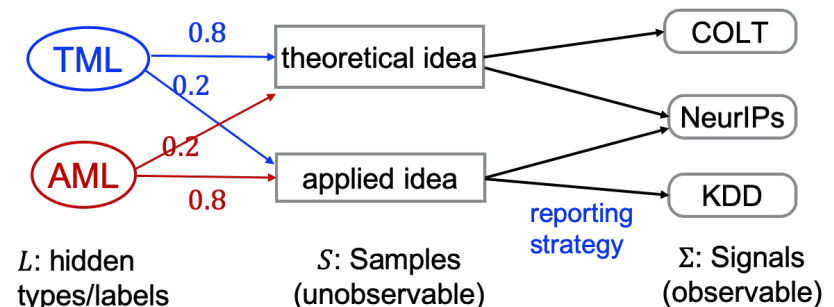
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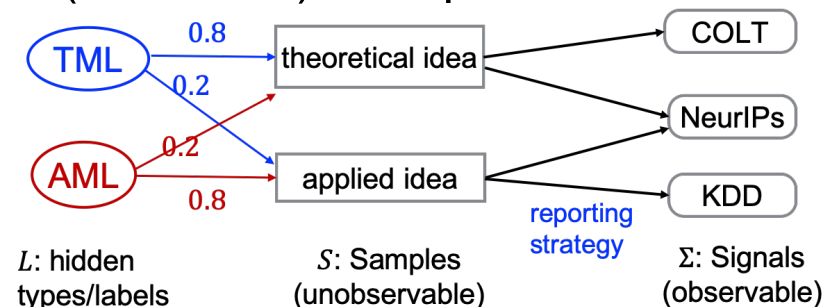
Model

- Two distribution types/labels: $l \in \{g, b\}$
 - g should be interpreted as “desired”, not necessarily good or bad
- $g, b \in \Delta(S)$ where S is the set of **samples**
- **Bipartite graph** $G = (S \cup \Sigma, E)$ captures feasible **signals** for each sample: $(s, \sigma) \in E$ iff σ is a valid signal for s
- g, b, G publicly known; S, Σ both discrete
- Distribution $l \in \{g, b\}$ generates T **samples**



Model

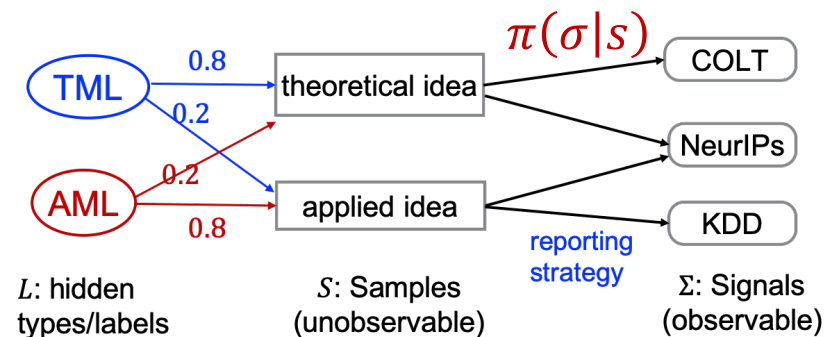
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- g, b, G publicly known; S, Σ both discrete
- Distribution $l \in \{g, b\}$ generates T **samples**
- A few special cases
 - Agent can hide samples, as in last lecture (captured by adding a “empty signal”)
 - Signal space may be the same as samples (i.e., $S = \Sigma$); G captures feasible “lies”



The Game

Agent's **reporting strategy** π transform T samples to a set R of T signals

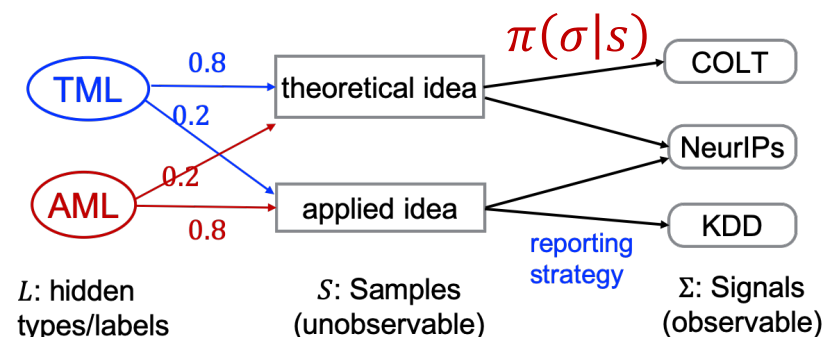
- A reporting strategy is a **signaling scheme**
 - Fully described by $\pi(\sigma|s)$ = prob of sending signal σ for sample s
 - $\sum_{\sigma} \pi(\sigma|s) = 1$ for all s



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- A reporting strategy is a **signaling scheme**
 - Fully described by $\pi(\sigma|s)$ = prob of sending signal σ for sample s
 - $\sum_{\sigma} \pi(\sigma|s) = 1$ for all s
- Given T samples, π generates T signals (possibly randomly) as an **agent report** $R \in \Sigma^T$
- A special case is deterministic reporting strategy



The Game

Agent's **reporting strategy** π transform T samples to a set R of T signals

➤ **Objective**: maximize probability of being accepted

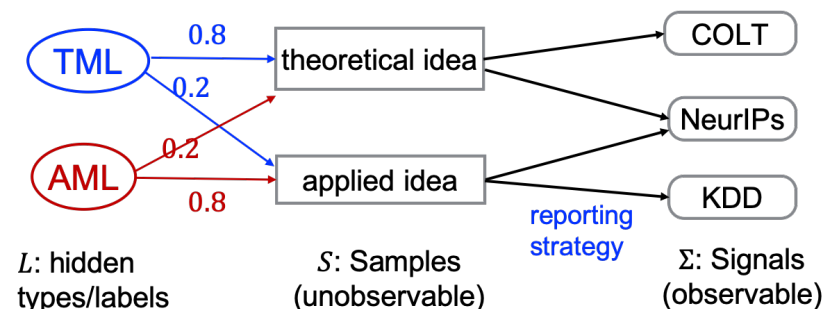
Principal's action $f: \Sigma^T \rightarrow [0,1]$ maps agent's report to an acceptance prob

➤ **Objective**: minimize prob of mistakes (i.e., reject g or accept b)

Remark:

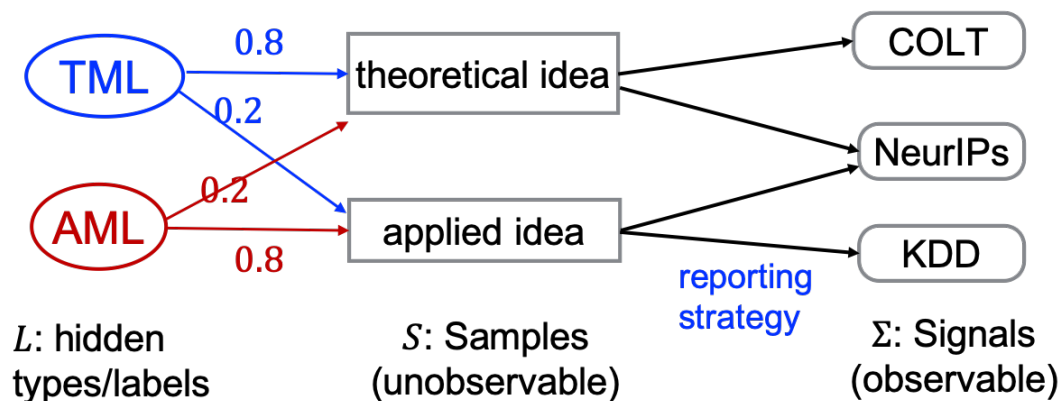
➤ Timeline: principal announces f first; agent then best responds

➤ Type g 's [b 's] incentive is **aligned with** [**opposite to**] principal



A Simpler Case

- Say $l \in \{g, b\}$ generates $T = \infty$ many samples
- Any reporting strategy π generates a distribution over Σ
 - $\Pr(\sigma) = \sum_{s \in S} \pi(\sigma|s) \cdot l(s) = \pi(\sigma|l)$ (slight abuse of notation)
 - $\pi(\sigma|l)$ is **linear** in variables $\pi(\sigma|s)$
- Intuitively, type g should make his π “far from” other’s distribution
 - Total variance (TV) distance turns out to be the right measure



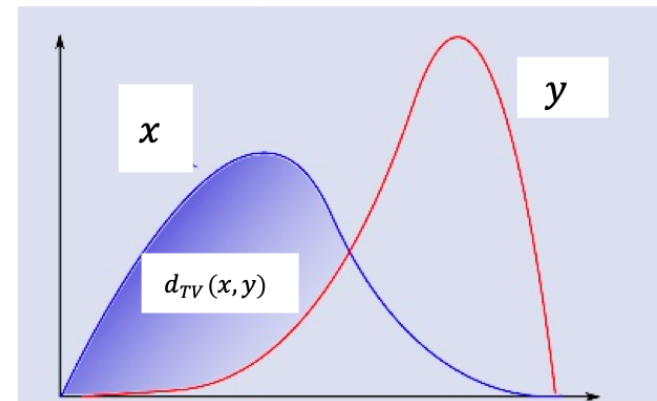
Total Variance Distance

➤ Discrete distribution x, y supported on Σ

- Let $x(A) = \sum_{\sigma \in A} x(\sigma) = \Pr_{\sigma \sim x}(\sigma \in A)$

$$\begin{aligned} d_{TV}(x, y) &= \max_A [x(A) - y(A)] \\ &= \sum_{\sigma: x(\sigma) > y(\sigma)} [x(\sigma) - y(\sigma)] \\ &= \underbrace{\frac{1}{2} \sum_{\sigma: x(\sigma) > y(\sigma)} [x(\sigma) - y(\sigma)]}_{\text{These two terms are equal}} + \underbrace{\frac{1}{2} \sum_{\sigma: y(\sigma) \geq x(\sigma)} [y(\sigma) - x(\sigma)]}_{\text{These two terms are equal}} \end{aligned}$$

These two terms are equal

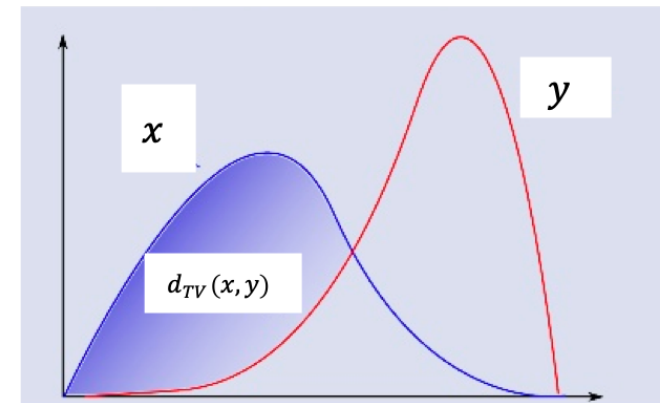


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How Can g Distinguish Himself from b ?

- Type g uses reporting strategy π (and b uses ϕ)
- Type g wants $\pi(\cdot | g)$ to be far from $\phi(\cdot | b)$ → What about type b ?
- This naturally motivates a zero-sum game between g, b

$$\max_{\pi} \min_{\phi} d_{TV} (\pi(\cdot | g) , \phi(\cdot | b)) = d_{DTV}(g, b)$$

Game value of this
zero-sum game

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Note $d_{DTV}(g, b) \geq 0$now, what happens if $d_{DTV}(g, b) > 0$?

- g has a strategy π^* such that $d_{TV}(\pi^*(\cdot | g), \phi(\cdot | b)) > 0$ for any ϕ
- Using π^* , g can distinguish himself from b with constant probability via $\Theta\left(\frac{1}{(d_{DTV}(g, b))^2}\right)$ samples
 - Recall: $\Theta\left(\frac{1}{\epsilon^2}\right)$ samples suffice to distinguish x, y with $d_{TV}(x, y) = \epsilon$
 - Principal only needs to check whether report R is drawn from $\pi^*(\cdot | g)$ or not

How Can g Distinguish Himself from b ?

- So $d_{DTV}(g, b) > 0$ is sufficient for distinguishing g from b
- It turns out that it is also necessary

Theorem:

1. If $d_{DTV}(g, b) = \epsilon > 0$, then there is a policy f that makes mistakes with probability δ when #samples $T \geq 2 \ln\left(\frac{1}{\delta}\right) / \epsilon^2$.
2. If $d_{DTV}(g, b) = 0$, then no policy f can separate g from b regardless how large is #samples T .

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Remarks:

- Prob of mistake δ can be made arbitrarily small with more samples
- We have shown the first part
- Second part is more difficult to prove, uses an elegant result for matching theory

But...Deciding Whether $d_{DTV}(g, b) > 0$ is Hard

Theorem: it is NP-hard to check whether $d_{DTV}(g, b) = 0$ or not.

➤ Recall $d_{DTV}(g, b) = \max_{\pi} \min_{\phi} d_{TV}(\pi(\cdot | g), \phi(\cdot | b))$

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Q: What goes wrong?

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Q: What goes wrong?

- We can only solve **normal-form** zero-sum games in poly time
- In that case, utility fnc is linear in both players' strategies
 - Can generalize to concave-convex utility fnc
 - But here, utility fnc is convex in both player's strategies

But...Deciding Whether $d_{DTV}(g, b) > 0$ is Hard

Theorem: it is NP-hard to check whether $d_{DTV}(g, b) = 0$ or not.

➤ Recall $d_{DTV}(g, b) = \max_{\pi} \min_{\phi} d_{TV}(\pi(\cdot | g), \phi(\cdot | b))$

Corollary: it is NP-hard to compute g 's best strategy π^* .

Proof:

- Will argue if we can compute π^* , then we can check $d_{DTV}(g, b) = 0$ or not
 - Thus computing π^* must be hard (actually “harder” than checking $d_{DTV}(g, b) = 0$)
- If we computed π^* , to compute $d_{DTV}(g, b)$, we only need to solve $\min_{\phi} d_{TV}(\pi^*(\cdot | g), \phi(\cdot | b))$ which is convex in ϕ
 - Minimize convex fnc can be done efficiently in poly time (well-known)
- First **example of reduction** in this class

Some Remarks

- Separability is determined by some “distance” between g, b
 - A generalization of TV distance to strategic setting
 - The principal’s policy is relatively simple
 - It is more of our own job to distinguish ourselves from others, rather than the employer’s
- The model can be generalized to many “good” (g_i) and “bad” (b_j) distributions
 - Principal wants to accept any g_i and reject any b_j
 - Separability is determined by $\min_{i,j} d_{DTV}(g_i, b_j)$
- The agent’s reporting strategy can even be adaptive
 - i.e., the π is different for different samples and may depend on past signals
 - Results do not change

Next Lecture will talk about how to utilize strategic manipulations to induce desirable social outcome

Thank You

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